Facebook F4 & Google Search & Virtual Machines

Hung-Wei Tseng

Recap: GFS v.s. WAS v.s. Facebook

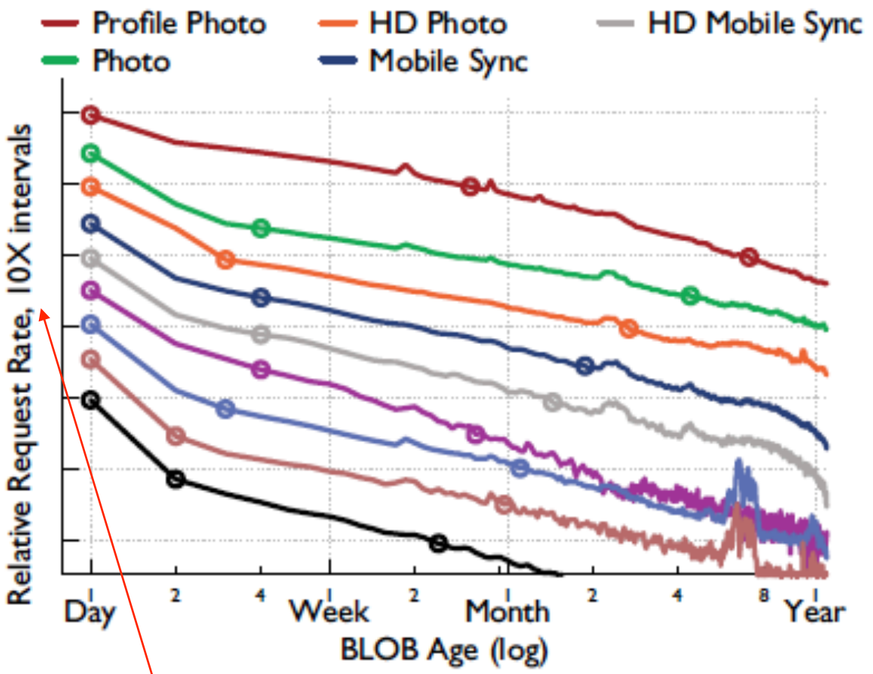
	GFS (OSDI 2003)	Facebook Haystack (OSDI 2010)	WAS (SOSP 2011)
File organizations	file chunk block	volume needle	stream extent record
System architecture	master chunkserver	directory haystack store	stream manager extent nodes
Data updates			append only updates
Consistency models	relaxed consistency		strong consistency
Data formats	files	photo/needle	multiple types of objects
Replications	intra-cluster replication	RIAD-6 & geo-replication	geo-replication
Usage of nodes	chunk server can perform both		separate computation and storage

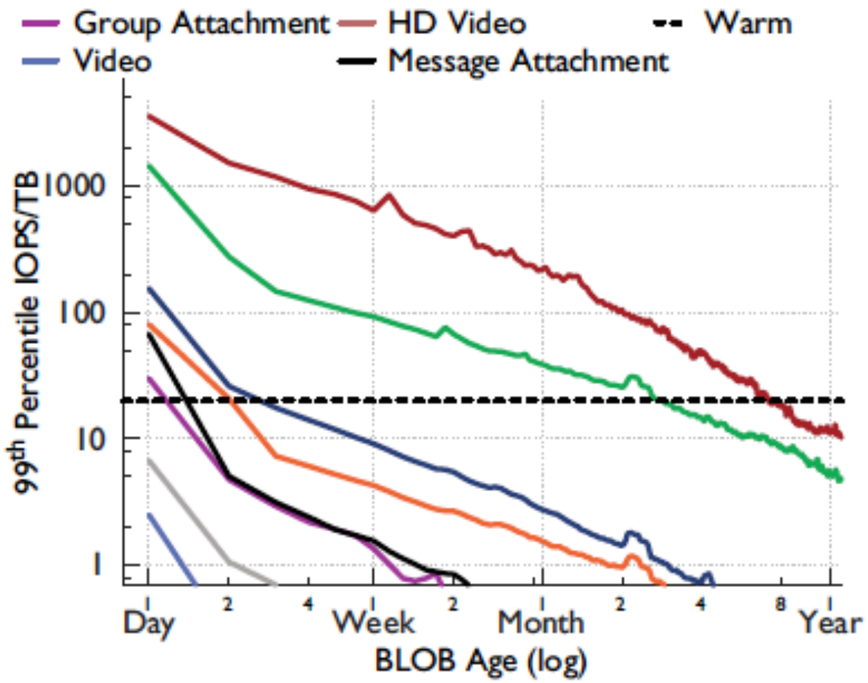
Recap: What kind of data is f4 optimized for?

- Regarding the type of data that f4 aims at, please identify how many the following statements is/are correct.
  - ① f4 is optimized for most frequently requested data in Facebook services
  - ② f4 is optimized for frequently created, deleted data
  - ③ f4 is optimized for reducing the access latency of long-term storage
  - ④ f4 is optimized for read-only data
  - ⑤ f4 is optimized for data that are not accessed very frequently

  A. 1
  B. 2
  C. 3
  D. 4
  E. 5

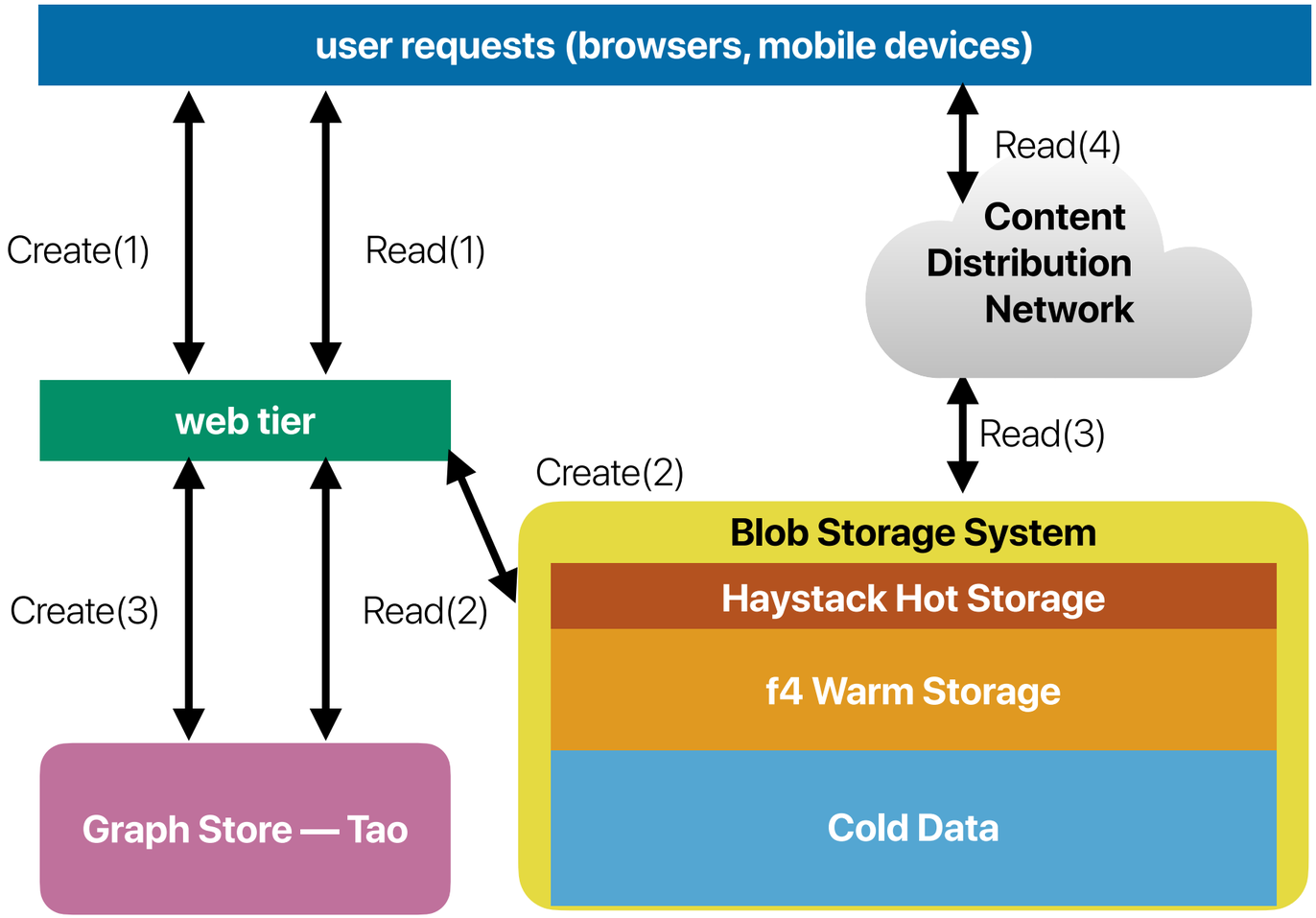
"Temperature" of data





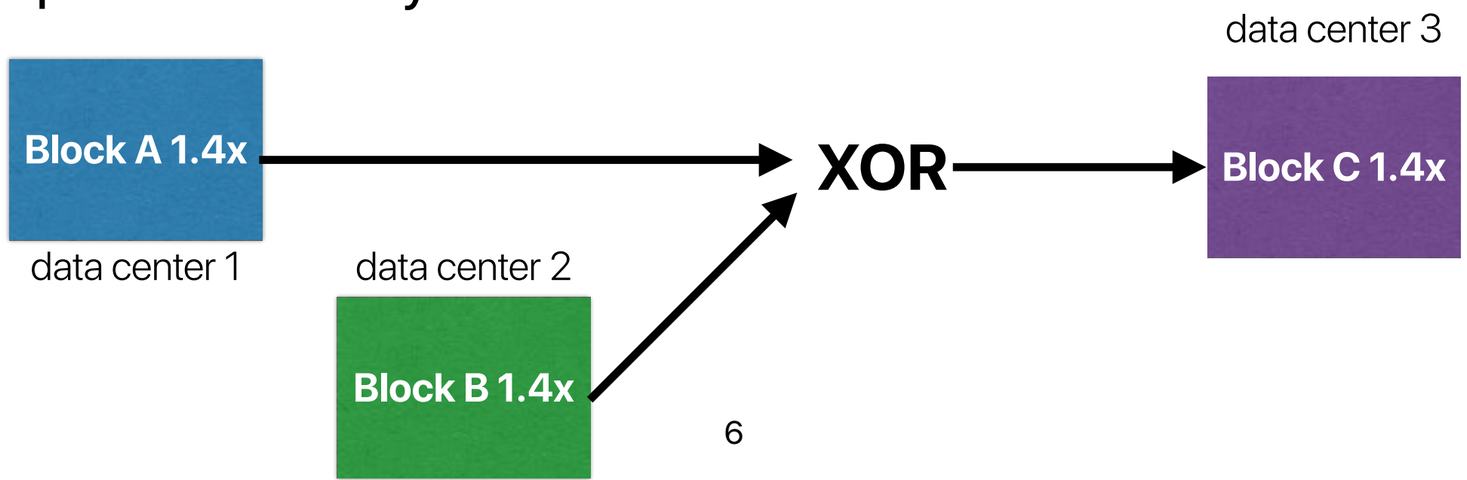
log scale — not encouraged to graph like this if you're writing a technical document or scientific paper

Recap: Facebook storage architecture



Recap: Storage efficiency

- Reed-Solomon erasure coding
  - Strips: 10GB data + 4GB parity — 1.4x space efficiency
  - One volume contains 10 strips
- XOR Geo-replication
  - Use XOR to reduce overhead further (e.g., Azure makes full copies)
  - Block A in DC1 + block B in DC2 -> parity block P in DC3
  - Any two blocks can be used to generate the third
  - 1.5x space efficiency
- 1.4*1.5 = 2.1x space efficiency in total



Current scoreboard

Red	Blue
21	20

Outline

- Facebook's F4 (cont.)
- Google Search
- Virtual machines

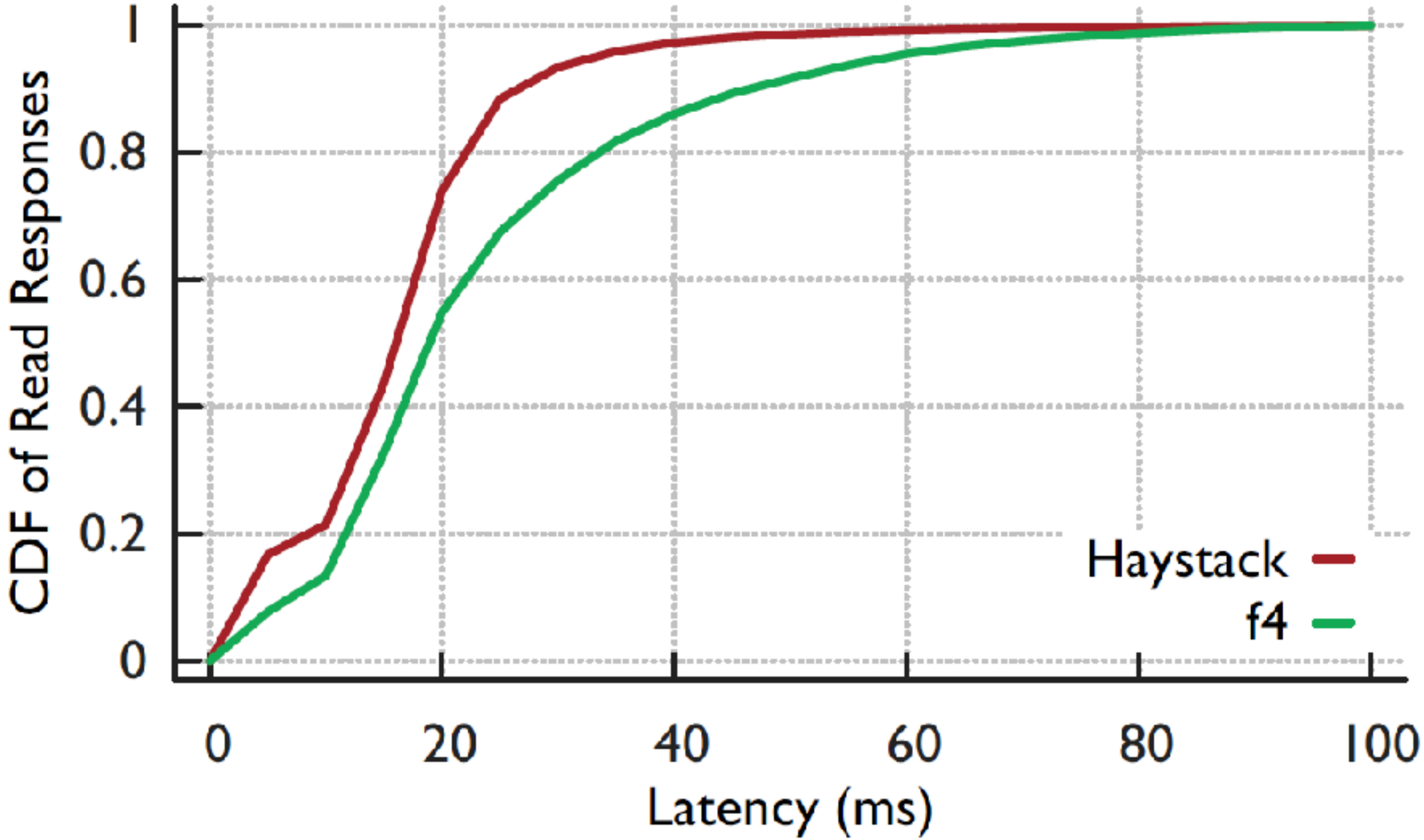
Fault tolerance

- 1%-2% HDD fail in a year
  - replicate data across multiple disks
  - Use erasure coding for storage efficiency
    - n blocks -> n + k blocks, can tolerate k simultaneous failures
    - higher cost for recovering data when there is a failure
- Host failures (periodically)
  - replicate coded blocks on different hosts
- Rack failures (multiple times/year)
  - replicate coded blocks on different racks
- Datacenter failures (rare, but catastrophic)
  - replicate blocks across data centers
  - use XOR to reduce overhead further (e.g., Azure makes full copies)
    - block A in DC1 + block B in DC2 -> parity block P in DC3
    - any two blocks can be used to generate the third
- Index files
  - use normal triple replication (tiny, little benefit in coding them)

What happens if fault occurs?

- Drive fails
  - Reconstruct blocks on another drive
  - Heavy disk, Network, CPU operation
  - one in background
- During failure, may need to reconstruct data online
  - rebuilder node reads BLOB from data + parity, reconstructs
  - only reads + reconstructs the BLOB (40KB), not the entire block (1GB)

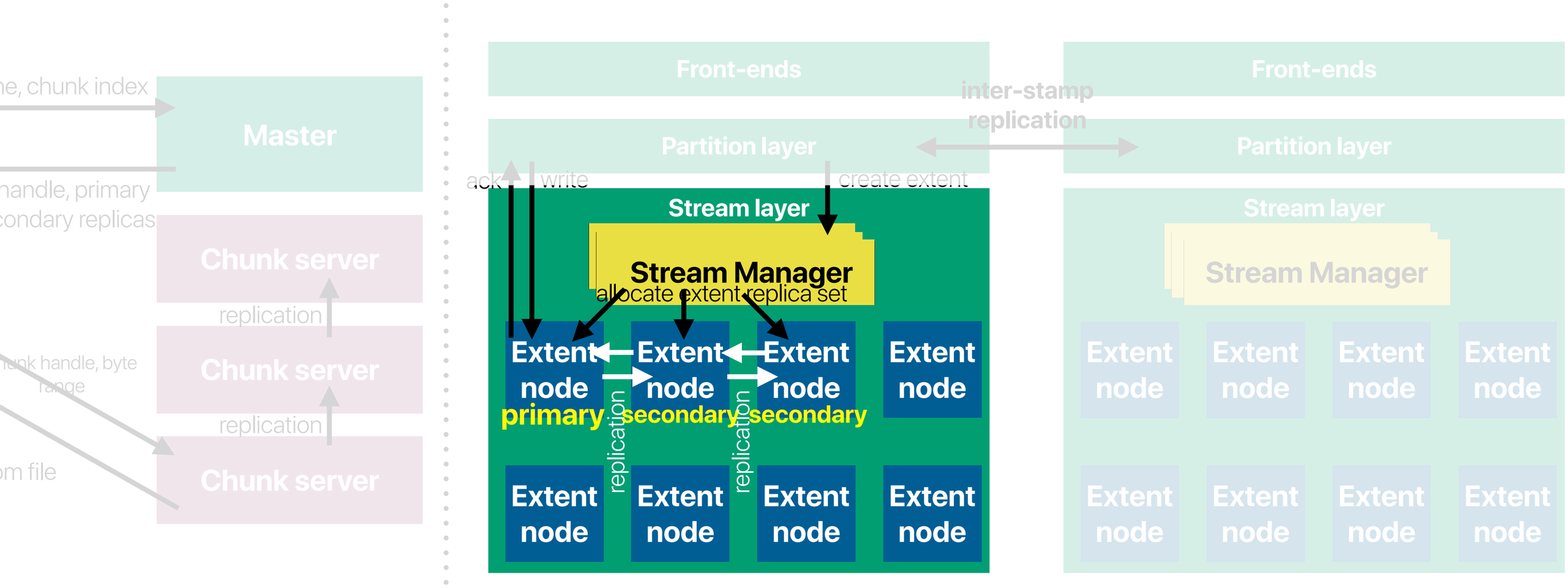
Performance of f4



Cells

- Each cell contains 14 racks of 15 hosts, each host contains 30 4TB H.D.Ds.
- A unit of acquisition, deployment
- Storage for a set of volumes
- Similar to the idea of stamps

Recap: GFS/WAS



Web search for a planet: The Google cluster architecture

Luiz Andre Barroso, Jeffery Dean ; Urs Holzle
Google

Poll close in 1:30

Google search architecture

- How many of the following fulfill the design agenda of the Google search architecture described in this paper?
  - ① Reduce the hardware cost by using commodity-class and unreliable PCs
  - ② Use RAID to provide efficiency and reliability
  - ③ Use replication for better request throughput and availability
  - ④ Optimize for the peak performance
  - A. 0
  - B. 1
  - C. 2
  - D. 3
  - E. 4

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  - ② Use RAID to provide efficiency and reliability
  - ③ Use replication for better request throughput and availability — replica, replica, replica
  - ④ Optimize for the peak performance — for performance per dollar

Also reliability and fault-tolerance

A. 0
B. 1
C. 2
D. 3
E. 4

- *Price/performance beats peak performance.* We purchase the CPU generation that currently gives the best performance per unit price, not the CPUs that give the best absolute performance.
- *Using commodity PCs reduces the cost of computation.* As a result, we can afford to use more computational resources per query, employ more expensive techniques in our ranking algorithm, or search a larger index of documents.

- *Software reliability.* We eschew fault-tolerant hardware features such as redundant power supplies, a redundant array of inexpensive disks (RAID), and high-quality components, instead focusing on tolerating failures in software.
- *Use replication for better request throughput and availability.* Because machines are inherently unreliable, we replicate each of our internal services across many machines. Because we already replicate services across multiple machines to obtain sufficient capacity, this type of fault tolerance almost comes for free.

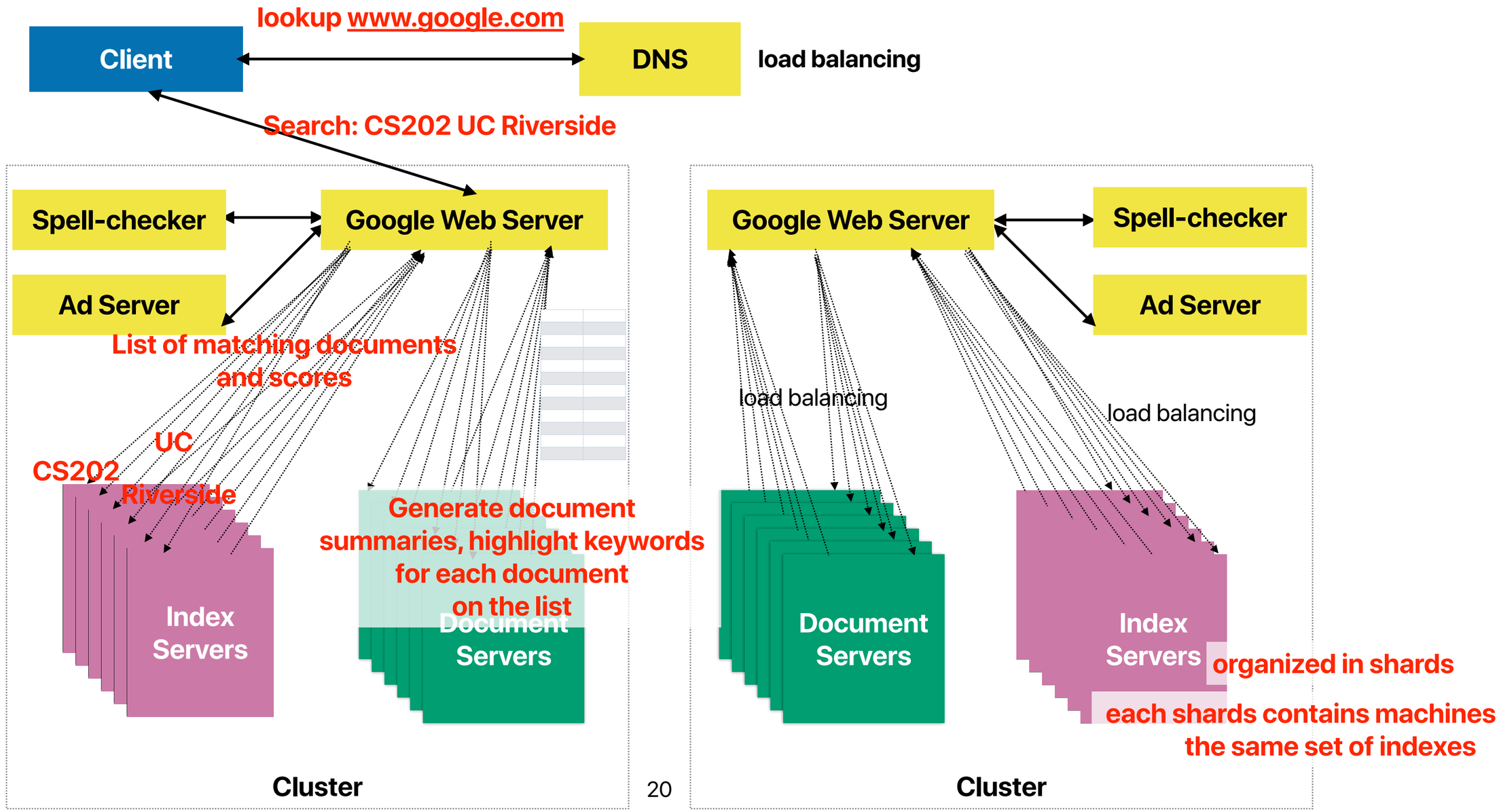
Why Google Search Architecture?

- The demand of performing search queries efficiently
  - Each query reads hundreds of MBs of data
  - Support the peak traffic would require expensive supercomputers or high-end servers
- We need a **cost-effective** approach to address this demand
  - Google search is compare against “AltaVista” search engine that uses DEC's high-performance alpha-based multiprocessor systems
  - AltaVista is later acquired by Yahoo! and you know the later story…

What Google proposes?

- Using commodity-class / unreliable PCs
- Provide reliability in software rather in hardware
- Target the best aggregate request throughput, not peak server response time

Google query-serving architecture



Replication is the key

- Scalability: simply add more replicas, the service capacity can improve
- Availability: even though one machine fails, another replica to take over

Poll close in 1:30

What kind of processors Google search needs

- If we are designing a processor just for Google search or similar type of applications, how many of the following targets/features would fulfill the demand?
  - ① Can execute many instructions from the same process/thread simultaneously
  - ② Can execute many processes/threads simultaneously
  - ③ Can predict branch outcome accurately
  - ④ Have very large cache capacity
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prediction logic. In essence, there isn't that much exploitable instruction-level parallelism (ILP) in the workload. Our measurements suggest that the level of aggressive out-of-order, speculative execution present in modern processors is already beyond the point of diminishing performance returns for such programs.

given how little ILP our application yields, and shorter pipelines would reduce or eliminate branch mispredict penalties. The avail-

end ones. Exploiting such abundant thread-level parallelism at the microarchitecture level appears equally promising. Both simultaneous multithreading (SMT) and chip multiprocessor (CMP) architectures target thread-level parallelism and should improve the performance of many of our servers. Some early

For such workloads, a memory system with a relatively modest sized L2 cache, short L2 cache and memory latencies, and longer (perhaps 128 byte) cache lines is likely to be the most effective.

Hardware

- Processor
  - Index search has little ILPs — doesn't need complex cores
  - Index search can be highly parallelized — processors with thread-level parallelism would be a good fit (e.g. Simultaneous Multithreading, SMT and Chip Multicprocessor, CMP)
  - Branch predictor matters
- Memory: Good spatial locality. Moderate cache size will suffice
- Storage: No SCSI, No RAID — not worth it
- Power: is an issue, but only $1,500/mo operating bill vs $7,700 capital expense

Will their architecture work for other things?

As mentioned earlier, our infrastructure consists of a massively large cluster of inexpensive desktop-class machines, as opposed to a smaller number of large-scale shared-memory machines. Large shared-memory machines are most useful when the computation-to-communication ratio is low; communication patterns or data partitioning are dynamic or hard to predict; or when total cost of ownership dwarfs hardware costs (due to management overhead and software licensing prices). In those situations they justify their high price tags.

At Google's scale, some limits of massive server parallelism do become apparent, such as the limited cooling capacity of commercial data centers and the less-than-optimal fit of current CPUs for throughput-oriented applications. Nevertheless, using inexpensive PCs to handle Google's large-scale computations has drastically increased the amount of computation we can afford to spend per query, thus helping to improve the Internet search experience of tens of millions of users. MICRO

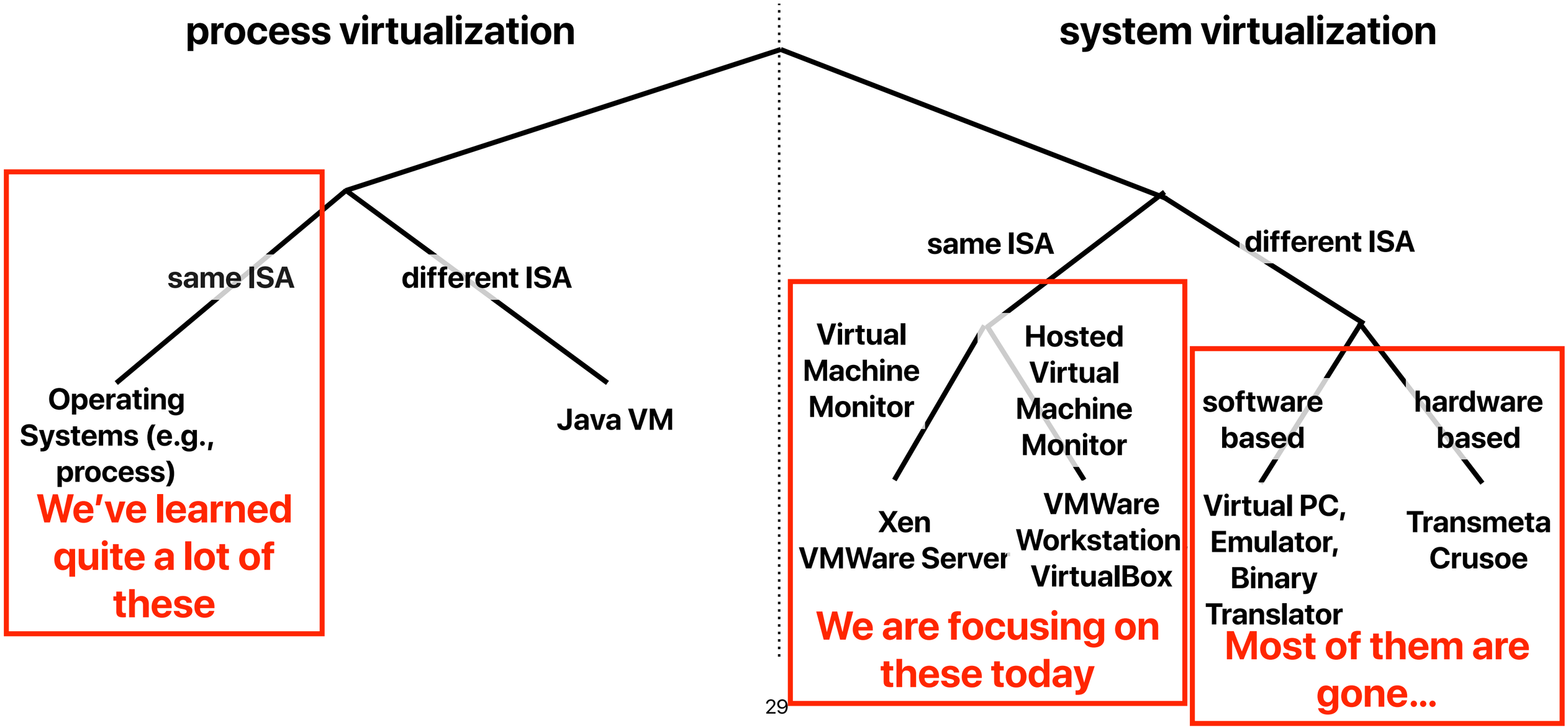
At first sight, it might appear that there are few applications that share Google's characteristics, because there are few services that require many thousands of servers and petabytes of storage. However, many applications share the essential traits that allow for a PC-based cluster architecture. As long as an application orientation focuses on the price/performance and can run on servers that have no private state (so servers can be replicated), it might benefit from using a similar architecture. Common examples include high-volume Web servers or application servers that are computationally intensive but essentially stateless. All of these applications have plenty of request-level parallelism, a characteristic exploitable by running individual requests on separate servers. In fact, larger Web sites already commonly use such architectures.

Metrics we care about data center design

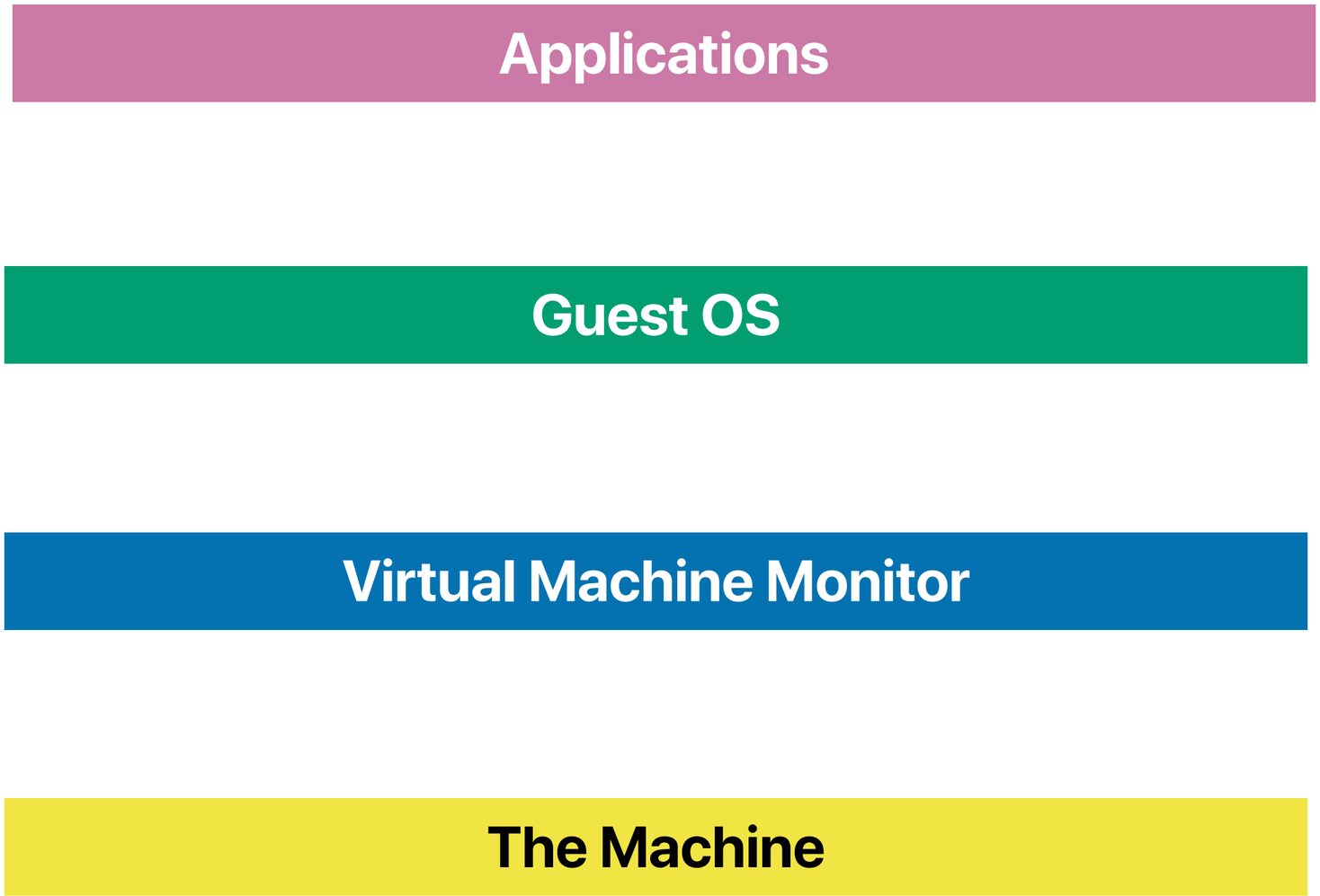
- Costs — machine architecture, distributed system architecture, replication strategies
- Power — machine architecture
- Energy — machine architecture
- Space-efficiency — erasure coding, replication, distributed
- Throughput — replication, distributed
- Reliability — replication

Virtual Machines

Taxonomy of virtualization



Virtual machine architecture



Three Laws of Robotics

- A robot may not injure a human being or, through inaction, allow a human being to come to harm.
- A robot must obey orders given it by human beings except where such orders would conflict with the First Law.
- A robot must protect its own existence as long as such protection does not conflict with the First or Second Law.

https://s3-ap-southeast-1.amazonaws.com/cloud-skcript/wp-content/uploads/2014/05/25090337/robots.jpg

Back to 1974...

Formal Requirements for Virtualizable Third Generation Architectures

Gerald J. Popek
University of California, Los Angeles
and
Robert P. Goldberg
Honeywell Information Systems and
Harvard University

Fidelity

Performance

Safety and isolation

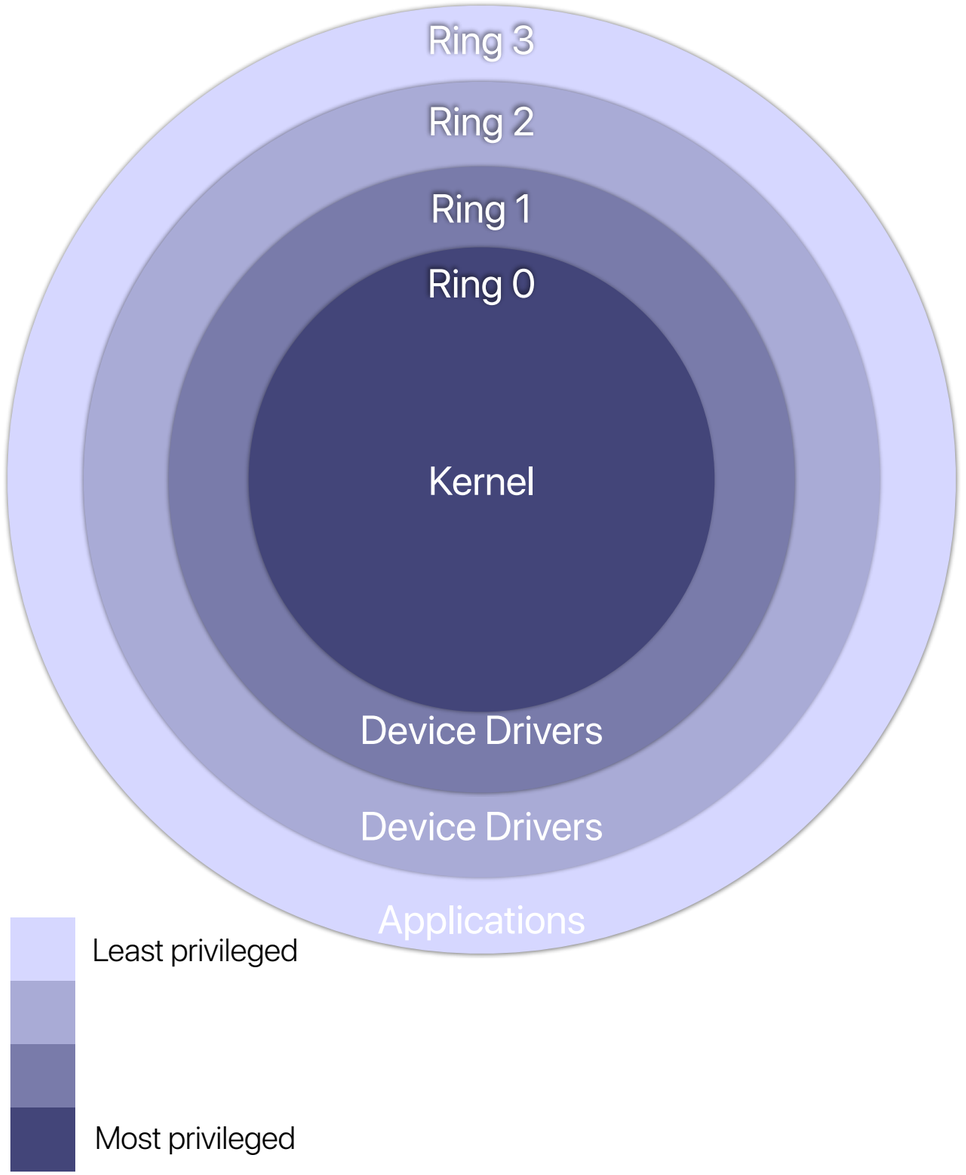
A virtual machine is taken to be an *efficient, isolated duplicate* of the real machine. We explain these notions through the idea of a *virtual machine monitor* (VMM). See Figure 1. As a piece of software a VMM has three essential characteristics. First, the VMM provides an environment for programs which is essentially identical with the original machine; second, programs run in this environment show at worst only minor decreases in speed; and last, the VMM is in complete control of system resources.

Recap: virtualization

However, we don't want everything to pass through this API!

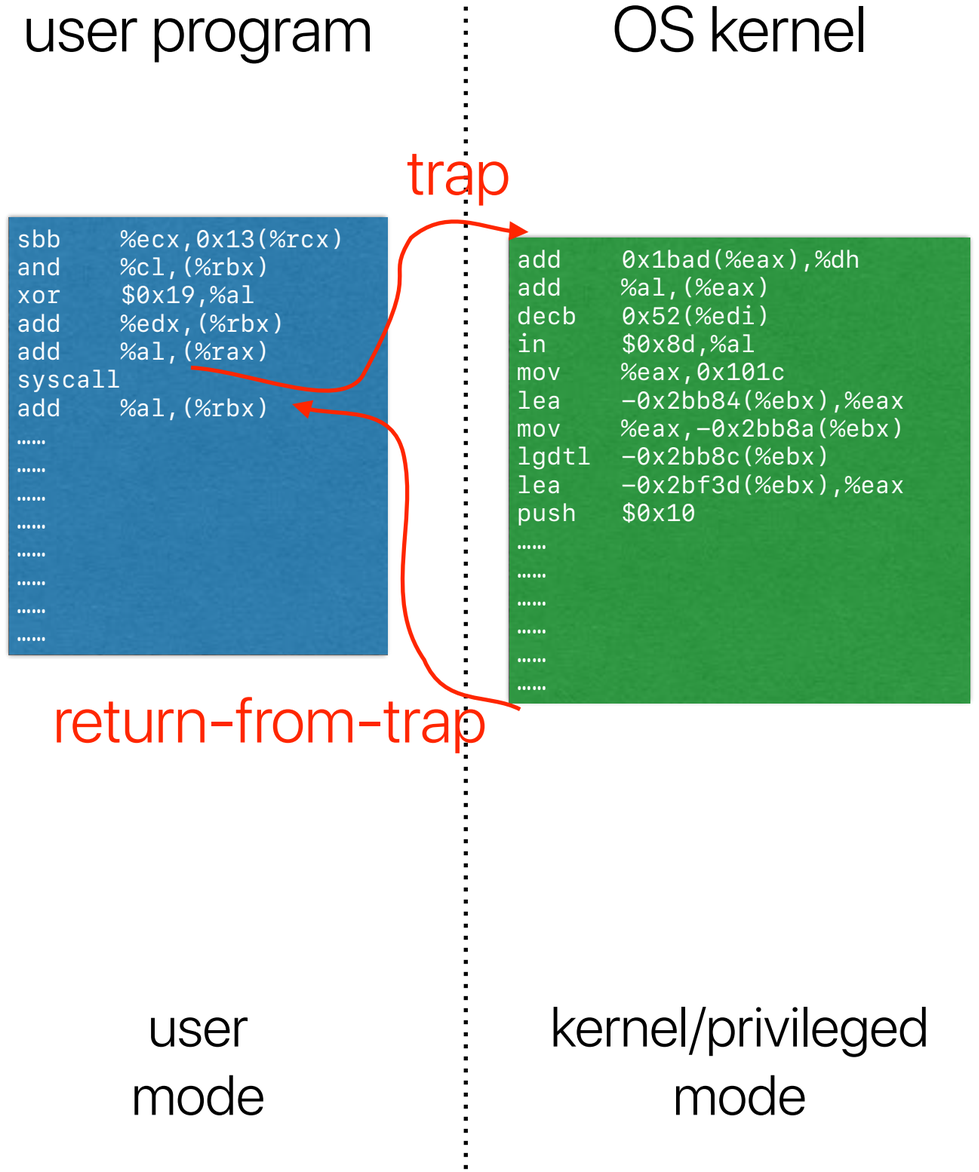
Recap: privileged instructions

- The processor provides **normal** instructions and **privileged** instructions
  - Normal instructions: ADD, SUB, MUL, and etc …
  - Privileged instructions: HLT, CLTS, LIDT, LMSW, SIDT, ARPL, and etc…
- The processor provides different modes
  - User processes can use normal instructions
  - Privileged instruction can only be used if the processor is in proper mode

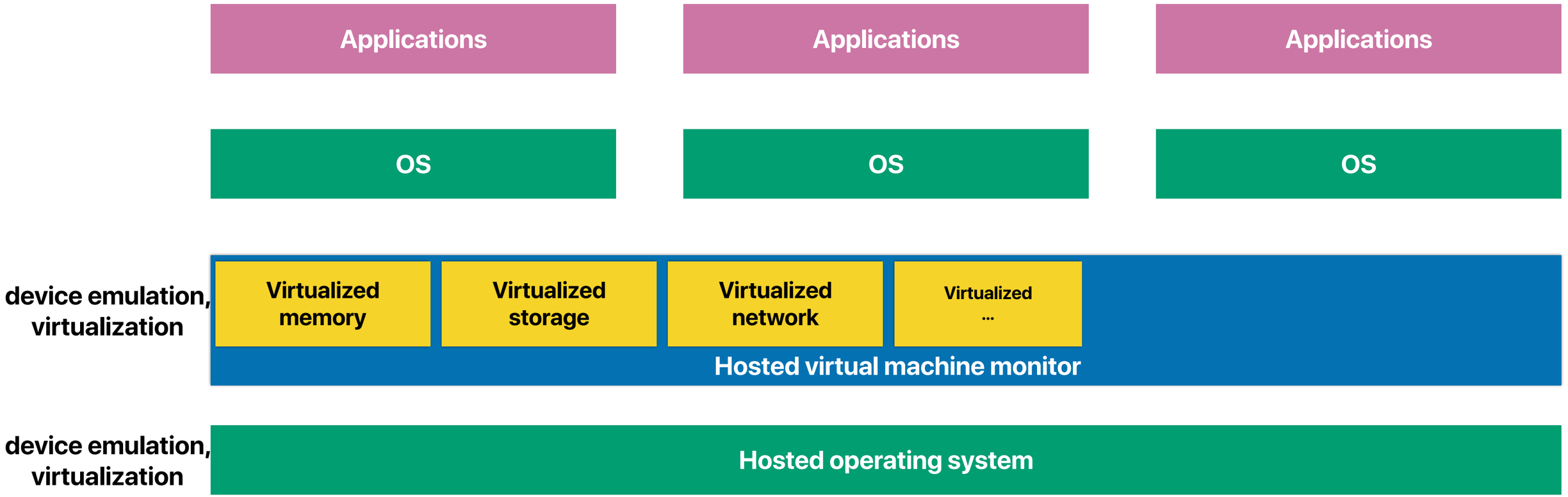


Recap: How applications can use privileged operations?

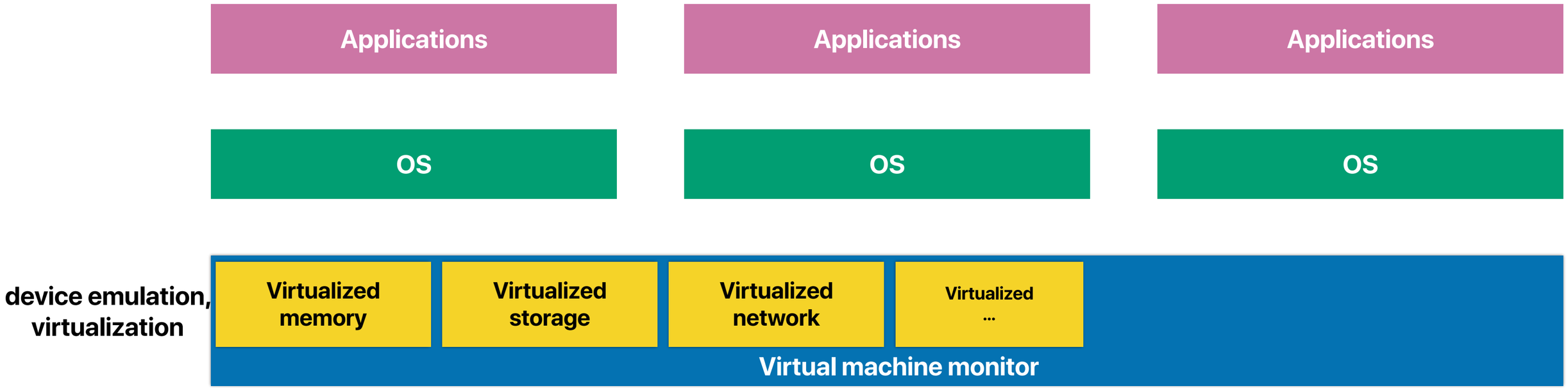
- Through the API: **System calls**
- Implemented in "trap" instructions
  - Raise an exception in the processor
  - The processor saves the exception PC and jumps to the corresponding exception handler in the OS kernel

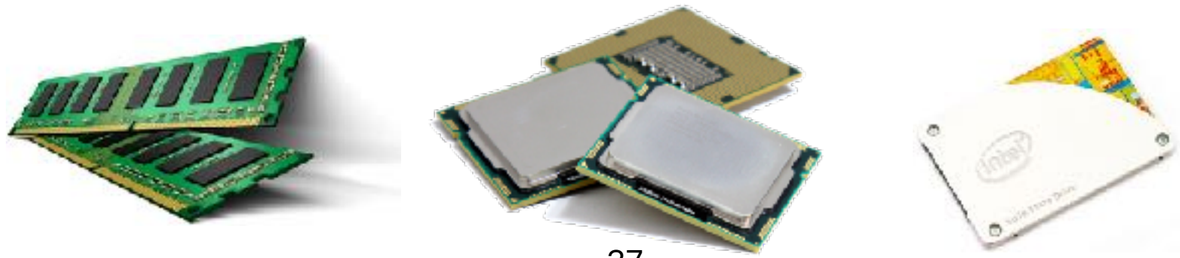


Hosted virtual machine



Virtual machine monitors on bare machines

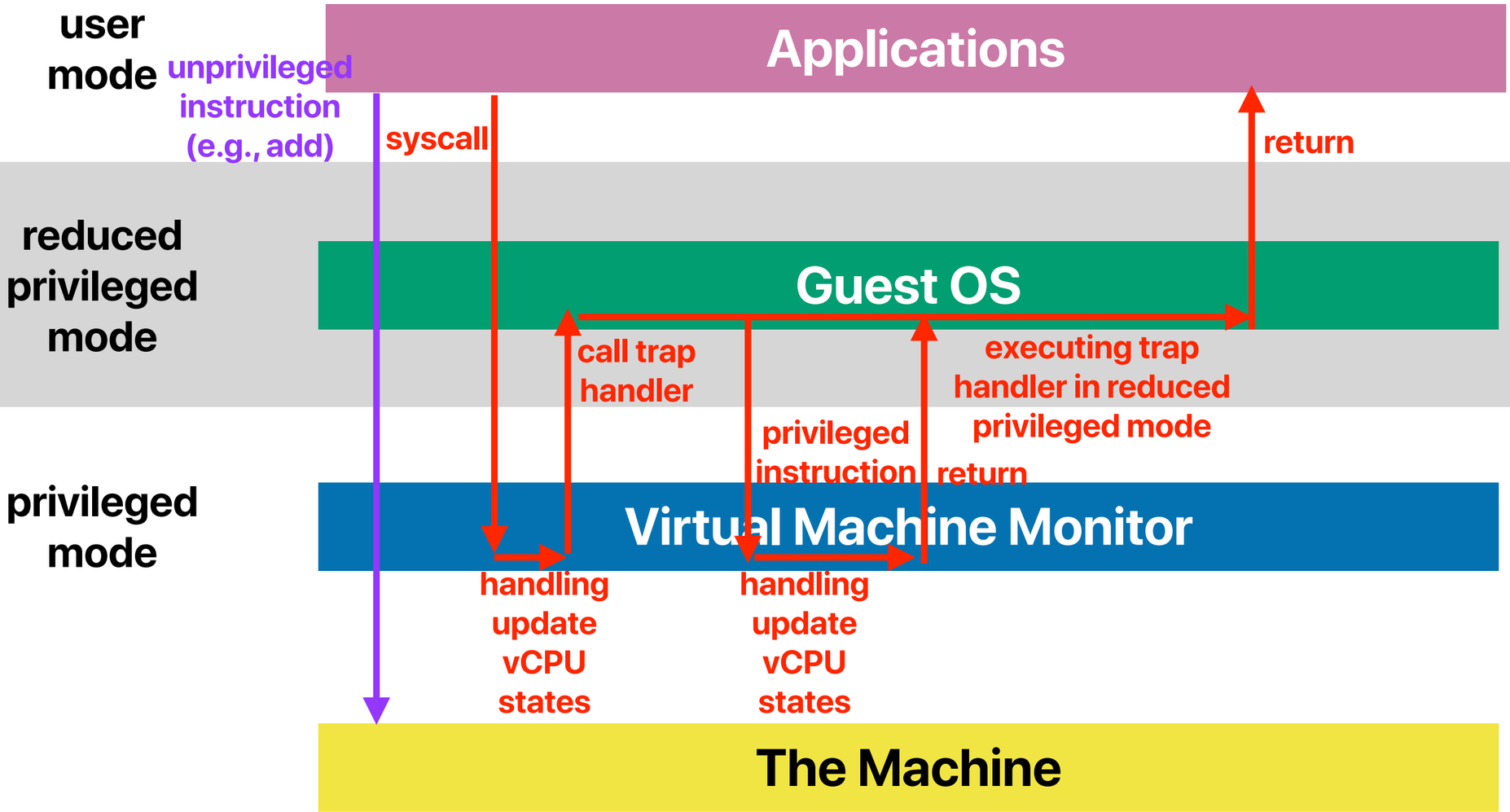




Three main ideas to classical VMs

- De-privileging
- Primary and shadow structures
- Tracing

CPU Virtualization: Trap-and-emulate



Announcement

- Group photo next lecture!
- Project revision
  - Allows you to revise your project with 20% of penalty on the unsatisfactory parts/test cases after the first-round of grading (firm deadline 3/11)
  - Say you got only 60% in the first-round, and you fixed everything before 3/11 — you can still get 60%+80%*40% = 92%
- iEVAL — count as an extra, full-credit reading quiz, due 3/11
- Final — contains two parts (each account for 50%)
  - Part 1: unlimited time between 3/12-3/17, open-ended questions
  - Part 2: 80 minute multiple choices/answers questions + two problem sets of comprehensive exam questions

Computer
Science &
Engineering

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